"Where the !?*! are the packets going?"

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Luca Sani

Senior R&D Software Engineer @Catchpoint







Traceroute

Traceroute is one of the most famous and long-lasting diagnostic tools in networking environment

First implementation by Van Jacobson in late 80s to answer the question:

"where the !?*! are the packets going" ?

Posted-Date: Tue, 20 Dec 88 05:13:28 PST Received-Date: Tue, 20 Dec 88 05:14:46 PST Received: from helios.ee.lbl.gov by venera.isi.edu (5.54/5.51) id AA25560; Tue, 20 Dec 88 05:14:46 PST Received: by helios.ee.lbl.gov (5.59/s2.2) id AA03127; Tue, 20 Dec 88 05:13:30 PST Message-Id: <8812201313.AA03127@helios.ee.lbl.gov> To: ietf@venera.isi.edu, end2end-interest@venera.isi.edu Subject: 4BSD routing diagnostic tool available for ftp Date: Tue, 20 Dec 88 05:13:28 PST From: Van Jacobson <van@helios.ee.lbl.gov> Content-Length: 2373 X-Lines: 46 Status: R0 After a frustrating week of trying to figure out "where the !?*!

After a frustrating week of trying to figure out "where the !?*! are the packets going?", I cobbled up a program to trace out the route to a host. It works by sending a udp packet with a ttl of one & listening for an icmp "time exceeded" message. If it gets one, it prints the source address from the icmp message, then bumps the ttl by one & etc. (As usual, I didn't come up with this clever idea -- I heard Steve Deering mention it at an end-to-end task force meeting.)

Traceroute implementations

- Many traceroute implementations have been created on different OSes
- Over the years it became one of the most used tools in the Internet measurement and topology discovery fields (multipath, de-aliasing, NAT traversal, ...)
 - Paris, Dublin, Pamplona traceroute...



Augustin, Brice, Timur Friedman, and Renata Teixeira. "Multipath tracing with Paris traceroute." Workshop on End-to-End Monitoring Techniques and Services. IEEE, 2007.

Linux traceroute

- We leverage Dmitry Butskoy's "<u>Linux tracerout</u>e"
 - Very fast
 - Open source
 - Easily extendible



- During the years we enhanced this traceroute to include new monitor capabilities
- We hope these enhancements can be useful to the community

Pietrasanta Traceroute

"A noble town since 1841 and a city of art" (and where our Italian office is located!)



Pietrasanta Traceroute

- QUIC traceroute
- ECN bleaching detection
- Work in Azure environment
- TCP "In Session"
- ... and many more

QUIC traceroute

QUIC

• QUIC is considered a transport layer protocol

- More than just "UDP"
- e.g., it is the transport layer of HTTP/3



QUIC assumes responsibility for the confidentiality and integrity protection of packets. For this it uses keys derived from a TLS handshake, but instead of carrying TLS records over QUIC (as with TCP), TLS handshake and alert messages are carried directly over the QUIC transport, which takes over the responsibilities of the TLS record layer.

<u>RFC9001</u> - Using TLS to Secure QUIC

QUIC support



- Packets sent are QUIC compliant, so the header is protected and the payload (frames) are encrypted
 - We leverage openssl3
- Nice "side effects"
 - Check whether the path filters QUIC
 - Determine if the destination supports QUIC
 - Check whether ECN is supported
 - Set IP-ECN in probes

QUIC traceroute

- Like "TCP half open"
- Do a QUIC handshake then closes the session (if opened)
 - Send QUIC "Initial" packet
 - Include a CRYPTO frame with TLS "Server Hello"
 - Intermediate hops will return ICMP TTL Exceeded
 - Destination may return
 - QUIC packet
 - ICMP port unreachable (still good, dest reached)
 - Nothing (timedout)
 - Close the session if it is the case
 - Send QUIC Initial packet including a CONNECTION_CLOSE frame

ECN bleaching detection

ECN mechanism

- ECN is a mechanism to signal that a packet experienced congestion (The Addition of Explicit Congestion Notification to IP <u>rfc3168</u>, **2001**)
- When a packet experiences congestion is marked instead of dropped
- The destination signals this event to the source, which in turn adjusts the rate

Network Working Group Request for Comments: 3168 Updates: 2474, 2401, 793 Obsoletes: 2481 Category: Standards Track K. Ramakrishnan TeraOptic Networks S. Floyd ACIRI D. Black EMC September 2001

The Addition of Explicit Congestion Notification (ECN) to IP

ECN marking

- Two bits into the IP header
- The source declares that a packet should be treated with ECN by setting the IP-ECN fields either to 01 or 10
- When congestion happens, instead of dropping the packet the router sets the IP-ECN fields to 11 (CE Congestion Experienced)



ECN feedback

- A destination that receives a packet with IP-ECN = CE should report to the source this event
- The source should then adjust the rate
- The report is done at transport/application layer
 - Example: in TCP, this event can be reported using a dedicated TCP flag (ECE ECN-Echo)

	++
	2 I S I F I
Header Length Reserved W C R C S S	5 Y I
R E G K H T	

ECN and L4S

- Recently, ECN mechanism got renewed attention due to L4S (Low Latency, Low Loss, and Scalable Throughput – <u>rfc9330</u>, **2023**)
- L4S requires an ECN feedback more accurate wrt the "classic" 2001 version



More accurate ECN feedback

TCP: More Accurate Explicit Congestion Notification (AccECN)
 Feedback in TCP (still a <u>draft</u>)

+++++	+++		1 +++++++	+
		А	C E U A P R S F	
Header Length	Reserved	Е	W C R C S S Y I	
			R E G K H T N N	
+++++	+++		<mark>.</mark> ++++++++	+

QUIC: Supported natively via <u>ECN counters</u> into the ACK frame

```
ECN Counts {
   ECT0 Count (i),
   ECT1 Count (i),
   ECN-CE Count (i),
}
```

ECN bleaching detection

- Intermediate hops can bleach/alter the value of ECN into the IP header (see for example: The Benefits of Using Explicit Congestion Notification (ECN) – <u>rfc8087</u>, 2017
- With Pietrasanta traceroute we can send probes with IP-ECN values different from zero and check hop by hop what was the IP-ECN value of the probe when it expired
- We can also check whether the destination transport layer (either TCP or QUIC) supports more accurate ECN feedbacks
 - TCP stack need to be patched
 - Not all QUIC implementations report ECN counters

ECN detection: Some examples

<pre>[bash]\$ sudo ./traceroute -nT -q 1ecn=1 -0 acc-ecn,info 95.2</pre>	28.44.181
traceroute to 95.228.44.181(95.228.44.181), 30 hops max, 60 byte	packets, ov
erall timeout not set	
1 172.21.82.1 <tos:1,dscp:0 ecn:1=""> (.234 ms</tos:1,dscp:0>	
2 64.79.149.27 <tos:1,dscp:<mark>),ECN:1> 1.374 ms</tos:1,dscp:<mark>	
3 64.79.139.17 <tos:1,dscp:<mark>),ECN:1> 1.297 ms</tos:1,dscp:<mark>	[hash 1\$ sudo /traceroute -nT -g 1ecn=1 -0 acc-ecn.info 81.236.63.162
4 66.209.72.25 <tos:1,dscp:0,ecn:1> 1.358 ms</tos:1,dscp:0,ecn:1>	f basis je state in the state of the stat
5 *	erall timeout not set
6 *	1 172 21 82 1 77051 DSCD+0 FCN+15 0 233 ms
7 4.68.39.58 <tos:1,dscp:0,ecn:1> 6.609 ms</tos:1,dscp:0,ecn:1>	2 = 64 - 70 - 440 - 27 - 7705 + 1 - 105 - 6 = CN + 1 - 1 - 276 mc
8 195.22.195.123 <tos:1,dscp:0,ecn:1> 160.604 ms</tos:1,dscp:0,ecn:1>	2 04./9.149.2/ <105.1,05CF.0,ECN.1> 1.2/0 ms
9 195.22.205.117 <tos:1,dscp:0,ecn:1> 173.535 ms</tos:1,dscp:0,ecn:1>	3 04./9.139.1/ <105.11,05.07.00,EUN:12 1.234 ms
	4 00.209.72.20 <105:1,05CP:0,ECN:12 1.271 ms
	5 60.209.04.124 <105:1, DSCP:0, ECN:12 1.115 ms
	6 62.115.32.150 < TOS:1, DSCP:0, ECN:1> 1.052 ms
	/ 62.115.132.119 <105:1,DSCP:0,ECN:1> 1.875 ms
17 95 228 44 181 <tos·1 dscp·0="" fcn·1=""> 170 007 ms</tos·1>	8 62.115.135.190 <105:1,DSCP:0,ECN:1> 6.789 ms
16 95.228.44.181 syn.ack.ece.cwr> 172.391 ms	9 62.115.137.38 <tos:1,dscp:0,ecn:1> 64.044 ms</tos:1,dscp:0,ecn:1>
Timedout: false	10 62.115.136.200 <tos:1,dscp:0,ecn:1> 69.195 ms</tos:1,dscp:0,ecn:1>
Duration: 1713.448 ms	11 80.91.254.90 <tos:1 145.761="" dscp:0,ecn:15="" ms<="" td=""></tos:1>
DestinationReached: true	12 62.115.139.172 <tos 1,dscp:0,ecn:1=""> 155.524 ms</tos>
	13 62.115.140.217 <tos 0,dscp:0,ecn:0=""> 150.248 ms</tos>
No bleaching destination	14 62.115.35.117 <tos:0,dscp:0,ecn:0> 150.434 ms</tos:0,dscp:0,ecn:0>
	15 81.228.89.186 <tos:0,dscp:0,ecn:0> 150.790 ms</tos:0,dscp:0,ecn:0>
supports AccECN over TCP	16 81.228.83.227 <tos:0,dscp:0,ecn:0> 150.816 ms</tos:0,dscp:0,ecn:0>
	17 90.228.166.164 <tos:0,dscp:0,ecn:0> 153.555 ms</tos:0,dscp:0,ecn:0>
	18 81.224.167.228 <tos:0,dscp:0,ecn:0> 153.135 ms</tos:0,dscp:0,ecn:0>
	19 *
	20 *
	21 81.236.63.162 <syn,ack> 150.907 ms</syn,ack>
Rleaching hannened	Timedout: false
Diedoning happened	Duration: 1522.420 ms
	DestinationReached: true

ECN bleaching research

- Run Pietrasanta traceroute from our vantage points deployed around the world to understand "how well" the network is ready to accommodate L4S
 - Where is the bleaching is happening?
 - Are there specific countries/ISPs/ASNs where it happens systematically?
- Stay tuned for more information!

City Overview									
							-		
	_		_		_				
Source	Ŧ	% Bleach	++	Avg RTT	\mathbf{v}	Avg Failing hop	$\overline{\mathbf{v}}$		
Jefferson		100		54		2			
Gilroy		100		62		3			
San Diego)	41		51		2			
Las Vegas	;	27		48		7			
Boston		25		50		4			
Phoenix		23		52		7			
New York		21		49		4			
Seattle		19		55		4			
Chicago		18		42		7			
Washingt	on (11		41		9			
Austin		11		63		13			
Denver		10		41		10			
Honolulu		9		93		7			
Kansas		9		35		7			
Walla Walla 9		56		10					



Work in Azure environment

Azure environment

- (Linux) VM with private IP
- Inbound ICMP packets are allowed

```
sudo traceroute -I google.com
traceroute to google.com (142.251.46.174), 30 hops max, 60 byte packets
1 * * *
2 * * *
3 * * *
4 * * *
5 * * *
6 * * *
7 * * *
8 * * *
9 * * *
10 * * *
11 nuq04s44-in-f14.1e100.net (142.251.46.174) 2.040 ms 2.050 ms 1.784 ms
```

- Intermediate hops are all *
- This happens for all traceroute protocols

Azure environment

- This happens because the source IP of the original probe encapsulated into the ICMP TTL Exceeded is left with the node public IP
- Thus, the ICMP reply is discarded by the kernel (not by traceroute)



Work in Azure environment

- We enhanced traceroute to work in "loose match mode"
- Open an additional raw ICMP socket to receive all ICMP packets and do the "kernel checks" at user level...
 - ... but do not check the source address of the encapsulated probe



TCP InSession

TCP "InSession"

- Classic TCP traceroute sends a different SYN for each hop
 - Different SYNs can take different paths
 - No consistency within a single traceroute
 - Many SYNs are sent per traceroute
 - Trigger firewall rules (SYN flood?)
- TCP InSession firstly opens a TCP session with the destination
- Then tracerouting is performed by sending 1-byte data packets within the session (with incremental TTL)

Checkout our blog for more information: <u>https://www.catchpoint.com/blog/traceroute-</u> <u>insession-catchpoints-effort-towards-a-more-reliable-network-diagnostic-tool</u>

And many more!

- Path MTU performance improvements
- Report ToS/DSCP hop by hop
- Report MSS when running in TCP mode
- Handle print in a separate thread (speed up)
- Overall timeout
- Compile and run on Alpine
- Avoid UDP standard filtering

Thank you!

- Feel free to check/use/ & contribute! <u>https://github.com/catchpoint/Networking.traceroute/</u> (GPL!)
- And come by to meet us!
 - Pietrasanta is a nice town on Tuscany seaside...





